

What System z Can Do That Intel-Based Systems Cannot

Chicago Area VM (and Linux) Enthusiasts (CAVMEN) 17 October 2013

What System z Can Do That Intel-Based Systems Can't

1. Transaction processing at scale

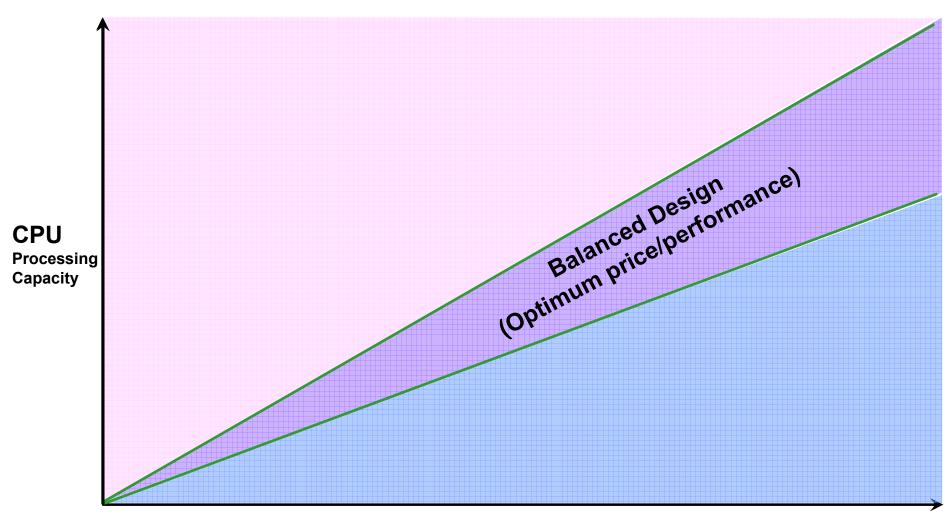


Servers Based on latest Intel technology (Sandy Bridge)



System z

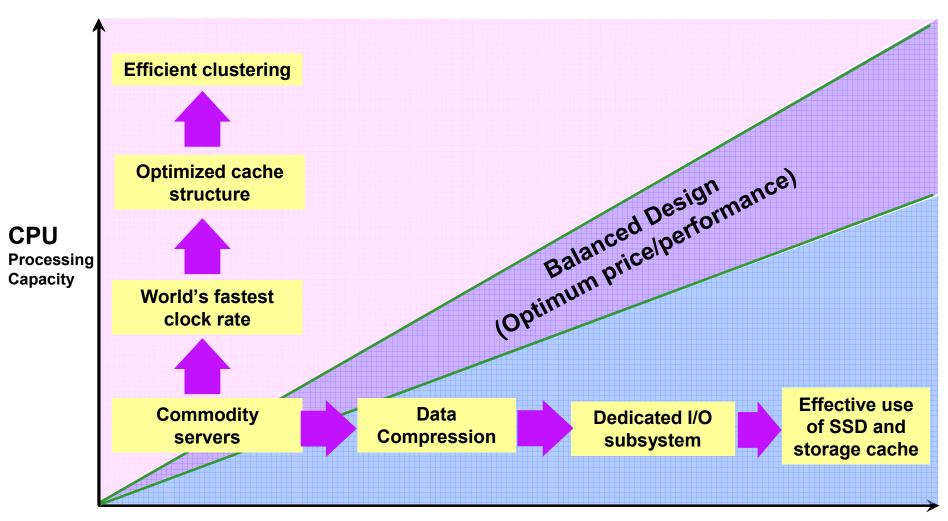
Transaction Processing At Scale Requires A Balance Of Capabilities



IOPS (Input Output Operations per Second)

What System z Can Do That Intel-based Systems Can't

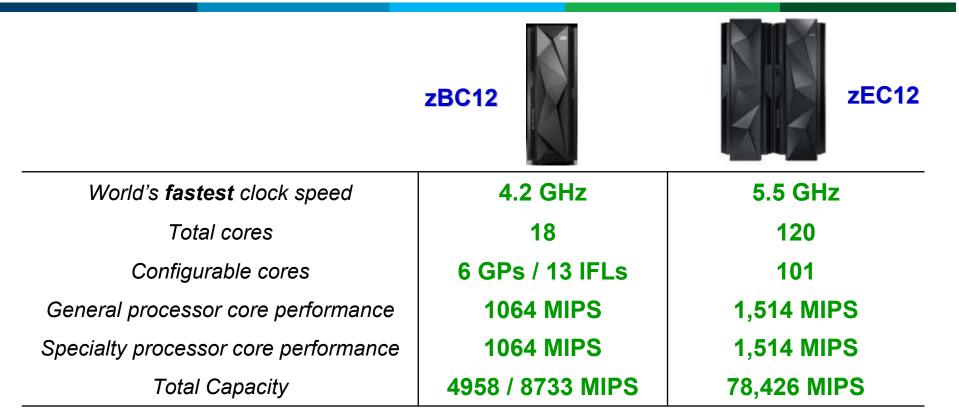
zEnterprise EC12 Design Is Unique And Optimized For Transaction Efficiency At Scale



IOPS (Input Output Operations per Second)

What System z Can Do That Intel-based Systems Can't

System z Delivers More Raw Processing Capacity Than Intel



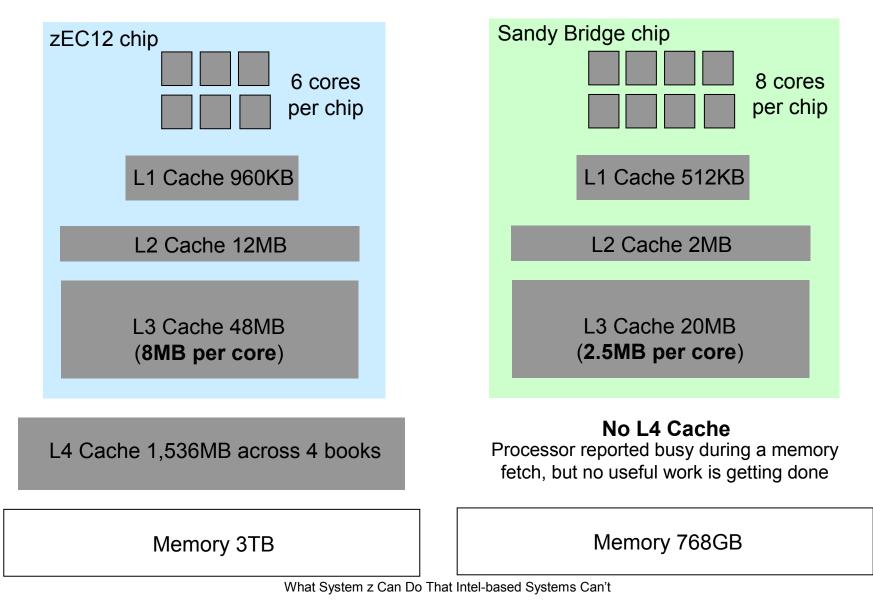


Maximum x86 clock speed = 3.4 GHz

Maximum x86 cores = 32

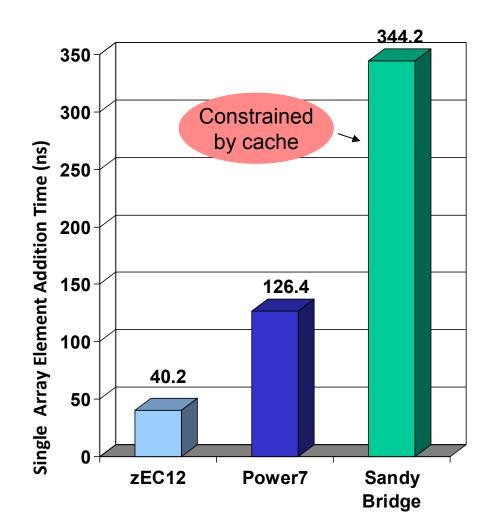
Intel Sandy Bridge

System z Has More Cache Than Intel To Support Multiple Concurrent Workloads



Intel Servers Slow Down Under Cache Intensive Workloads

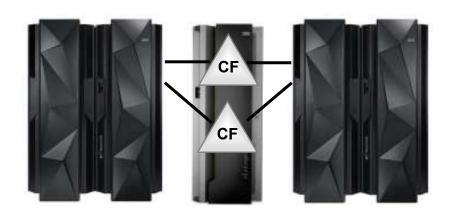
- Multiple concurrent processes introduces cache contention
 - Example: 5 processes each with 70MB working set size
- Intel workloads significantly slowed due to cache contention
- System z with z/OS showed results 8X faster than Intel system



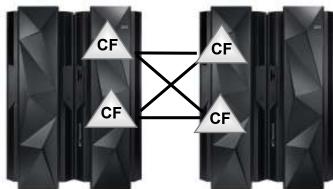
Parallel Sysplex Enables System z To Scale To Capacities Far Beyond What Intel Can



Parallel sysplex clustering delivers highest availability



Single System Sysplex



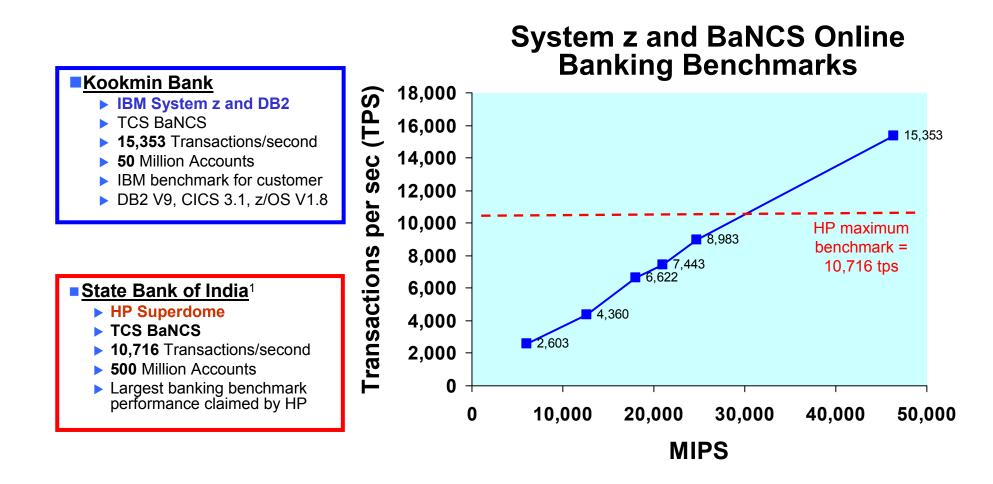
Cross Connected Servers with internal Coupling Facilities

External Coupling Facility (Can be different class server)

Potentially 2.5 million MIPS per 32-way cluster

Supports rolling software updates via automatic sysplex failover

Real-World Benchmarks Show System z Runs Bigger Workloads Than Intel



¹ Source: http://www.enterprisenetworksandservers.com/monthly/art.php?2976 and *InfoSizing FNS BANCS Scalability on IBM System z – Report Date: September 20, 2006;* Clement Report; http://h20195.www2.hp.com/v2/GetPDF.aspx/4AA1-4027ENW.pdf Feb 2010

Sustains High CPU Utilizations With No Service Degradation

16.00

16.15

16.30

16.45

17.00

99.60

97.40

92.40

90.50

94.20

06/10

06/10

06/10

06/10

06/10

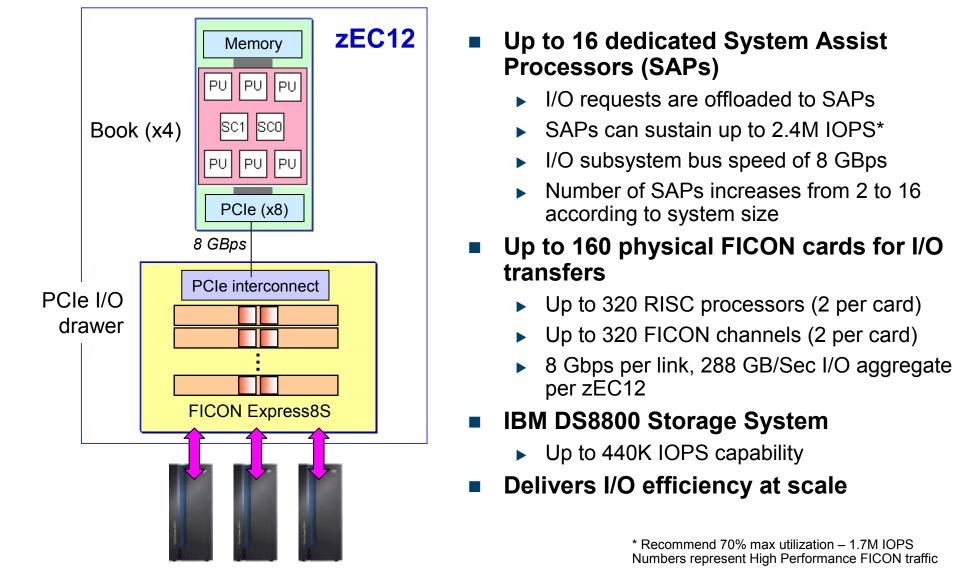
Recent Daily Samples Indicating Stress Extended Periods above 97%

Monday June 10, 2013

8:15 to Noon			Noon to 5PM		
06/10	8.15	91.90	06/10	12.15	97.00
06/10	9.15	91.00	06/10	12.30	97.00
06/10	9.30	95.30	06/10	12.45	90.40
06/10	9.45	99.10	06/10	13.00	91.20
06/10	10.00	99.90	06/10	13.15	92.90
06/10	10.15	98.80	06/10	13.30	95.70
06/10	10.30	99.80	06/10	13.45	96.00
06/10	10.45	99.80	06/10	14.00	99.40
06/10	11.00	100.10	06/10	14.15	100.00
06/10	11.15	100.10	06/10	14.30	99.10
06/10	11.30	99.40	06/10	14.45	98.10
06/10	11.45	97.60	06/10	15.00	97.60
06/10	12.00	99.70	06/10	15.15	99.60
	1410050304218		06/10	15.30	99.80
			06/10	15.45	99.00

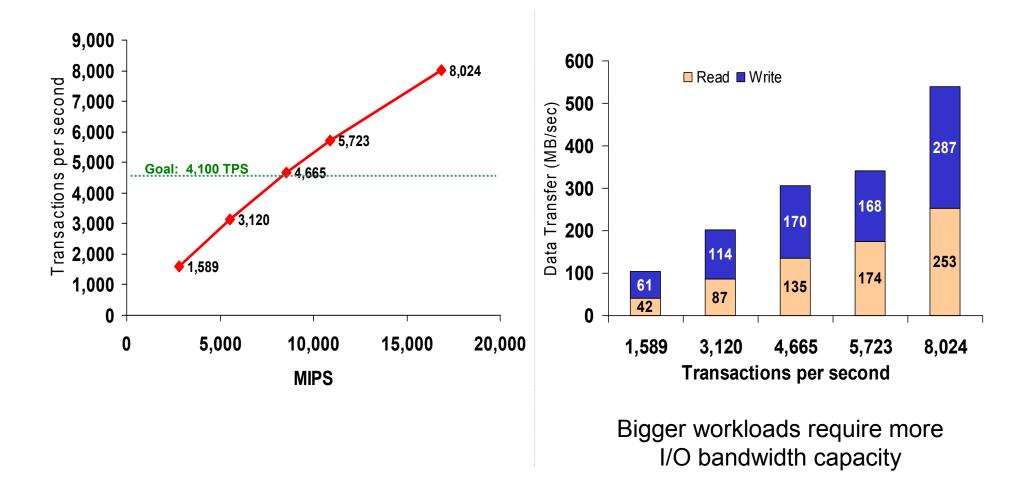
Sustained Peaks	Number of Days	Dates
100%	2	4,10
99%+	6	3,4,6,10,19, 25
97%+	9	3,4,6,10,11, 19,24,25,26

zEC12 Has A Dedicated I/O Subsystem Which Can Deliver 1.7M I/O Operations Per Second



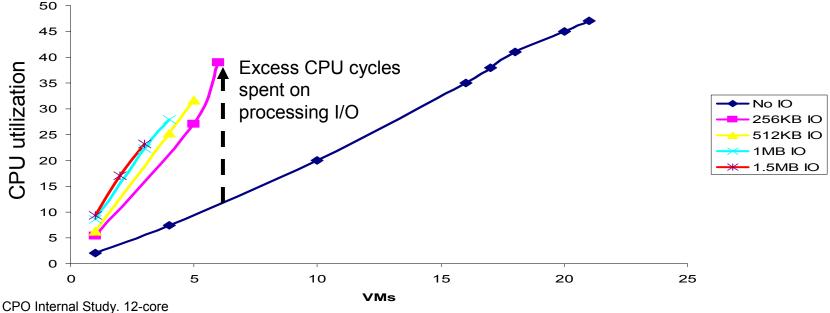
I/O Bandwidth Is Important For Critical Data Workloads

Bank of China System z Benchmark required huge I/O bandwidth capacity



System z's Dedicated I/O Subsystem Delivers More I/O Processing Capacity Than Intel

- Intel's performance degrades as I/O demand increases
 - No dedicated I/O subsystem
- Test case scenario: Run multiple virtual machines on x86 server
 - Each virtual machine has an average I/O rate
 - x86 processor utilization is consumed as I/O rate increases

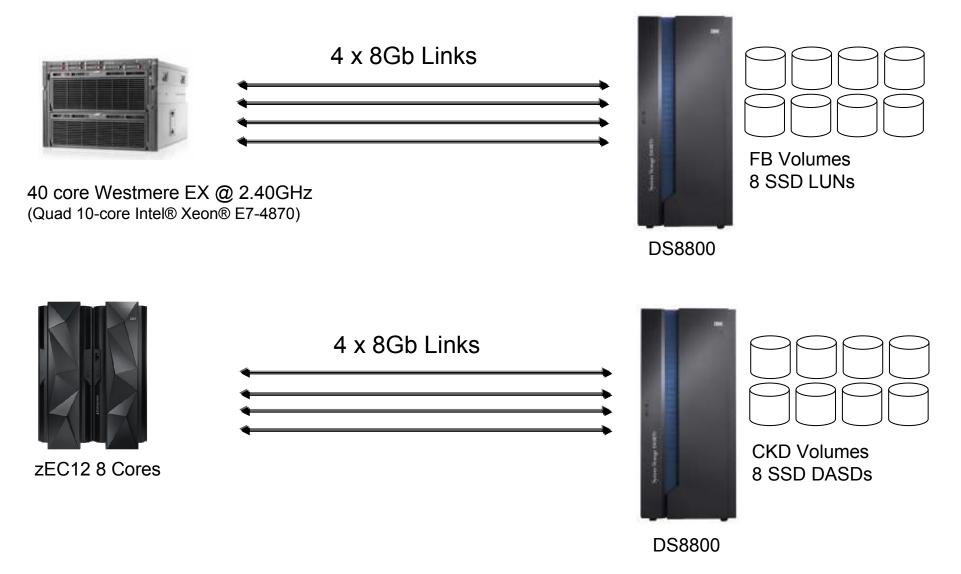


Intel CPU As IO Load Increases

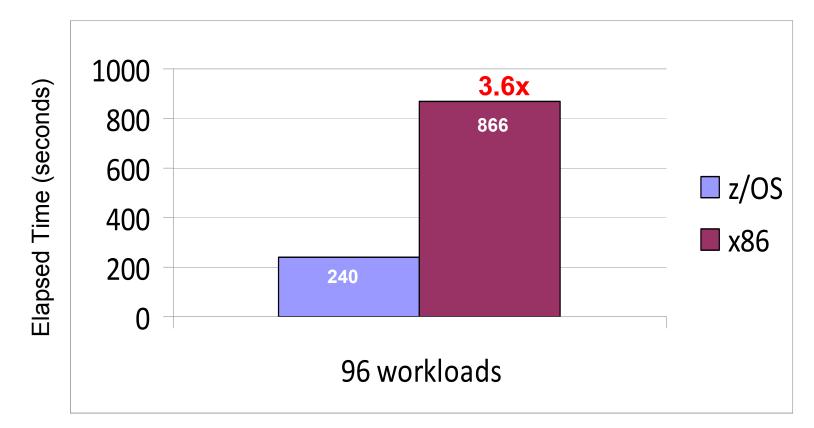
Source: CPO Internal Study. 12-core Westmere EP with KVM. FB at 22 tps with varying IO per transaction.

What System z Can Do That Intel-based Systems Can't

Test Configuration

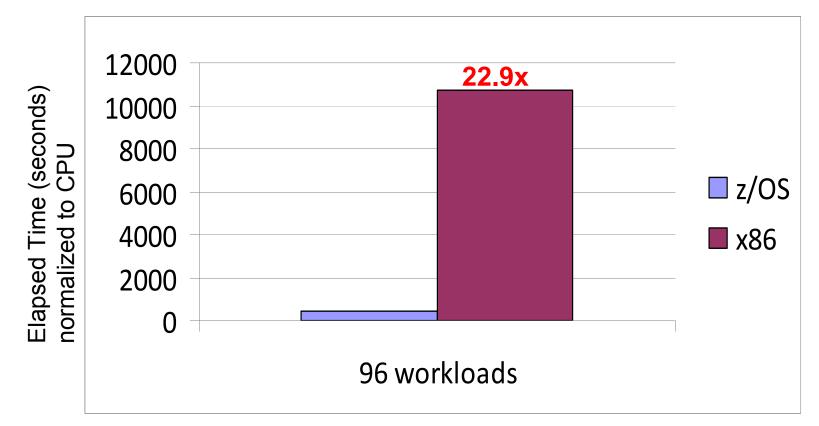


Elapsed Time Comparison



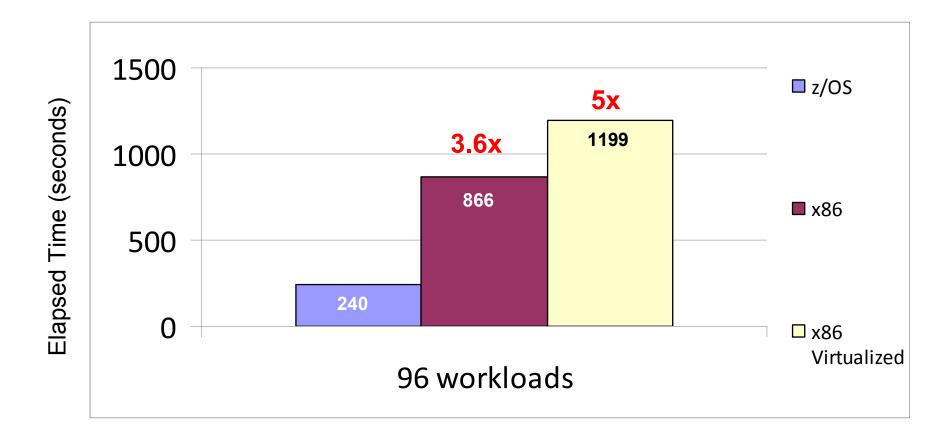
Important note: 96-wide was the max point to which the test cases could be run before the backend storage available for this study became a bottleneck. System z channel utilization is at 27% and the IOPs are only at 2.5%, leaving z considerable room to grow

Elapsed Time Normalized to CPU



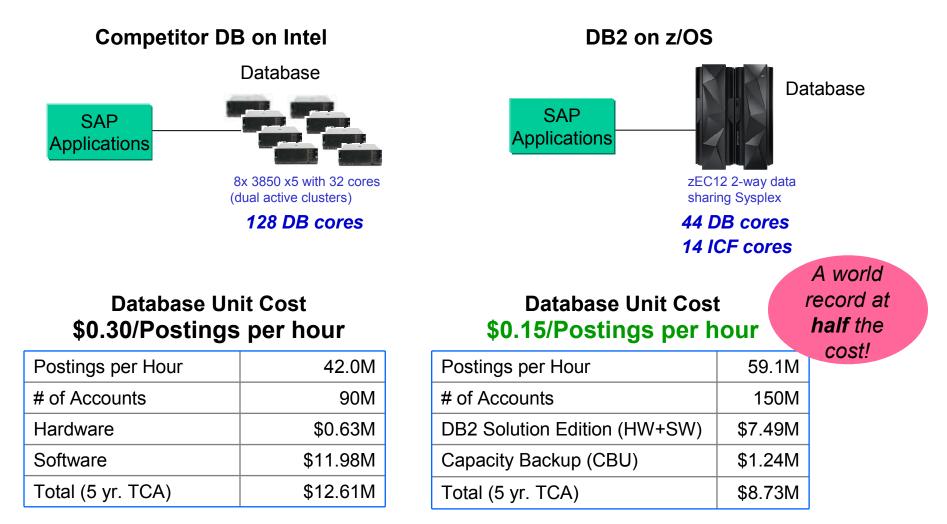
Important note: 96-wide was the max point to which the test cases could be run before the backend storage available for this study became a bottleneck. System z channel utilization is at 27% and the IOPs are only at 2.5%, leaving z considerable room to grow

Compare Impact of Virtualization



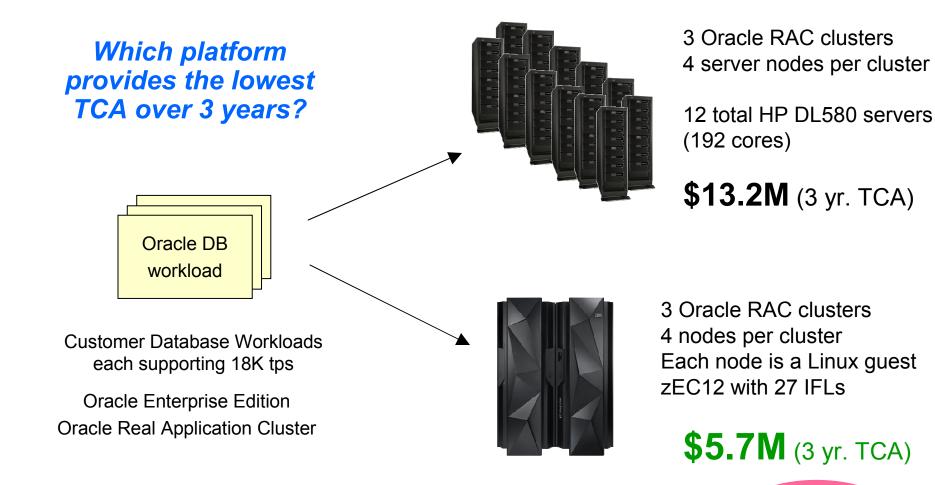
Important note: 96-wide was the max point to which the test cases could be run before the backend storage available for this study became a bottleneck. System z channel utilization is at 27% and the IOPs are only at 2.5%, leaving z considerable room to grow

z/OS Database Workloads Benefit From Higher I/O Bandwidth



Cost of platform infrastructure for comparative transaction production. Cost of packaged application software not included. List prices used.

Consolidated Oracle DB Workloads Benefit From Linux On System z's I/O Bandwidth



TCA includes hardware, software, maintenance, support and subscription. Workload Equivalence derived from a proof-of-concept study conducted at a large Cooperative Bank.

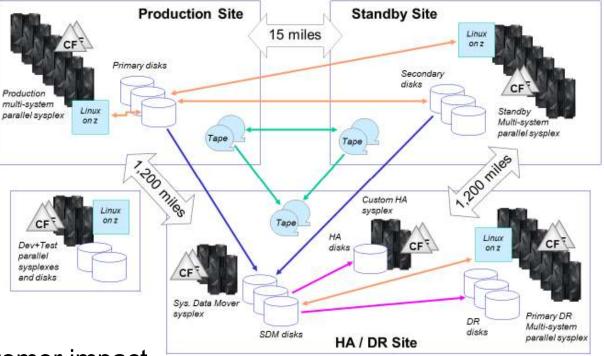
What System z Can Do That Intel-based Systems Can't

Half the

cost

Intel Does Not Have The Physical Capacity For State-of-the-Art Systems Of This Magnitude

- 1B CICS trans/day
- 4,000 IMS trans/sec
- 14M ACH transactions in 2.5 hours
 - 6-way sysplex
 - 30ms response
 - 216 CPU's at primary site
 - 200K MIPS



- Zero outages, zero customer impact
- Linux is Active-Active in the two data centers, with zero downtime
 - 15% Linux, growing at 30%
- *"Crazy about security overall, and the z system has a fortress around it"*

What System z Can Do That Intel-Based Systems Can't

1. Transaction processing at scale

2. Perfect Workload Management





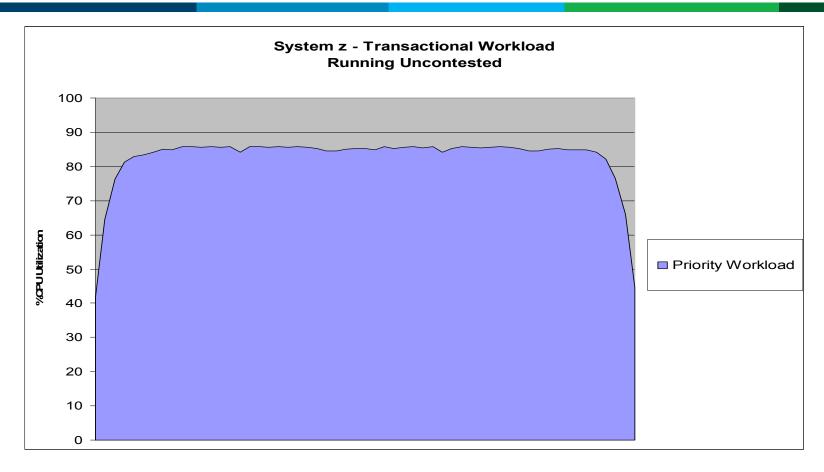
Servers Based on latest Intel technology (Sandy Bridge)

System z

Workload Management

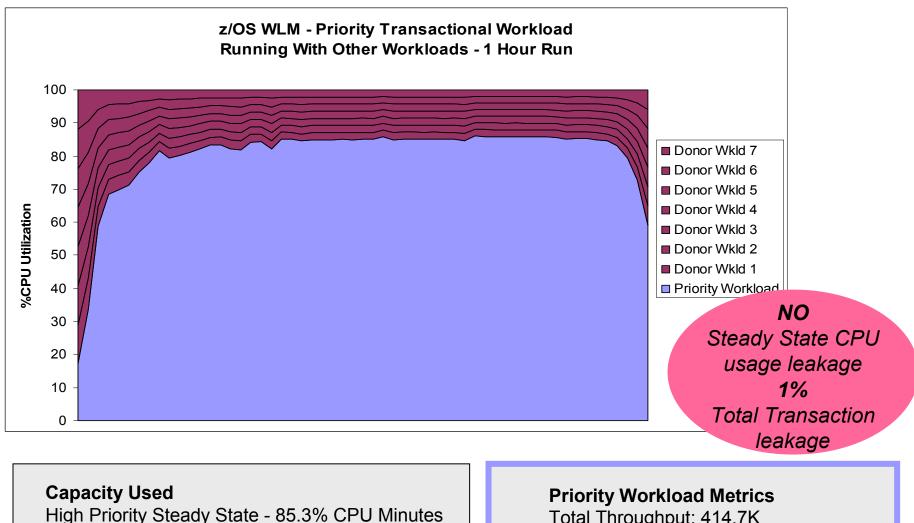
- Hosting platforms must be able to support high priority and low priority workloads together when sharing resources
 - Enables maximum utilization of the hosting platform
- Particularly relevant in a Private cloud environment
 - Multiple tenants with different priorities
- Desired behavior when mixing workloads
 - Low priority workloads "give up" resources to high priority workloads when required, soak up unused resources when available
 - High priority workload performance must not degrade

Priority Transactional Workload With Constant Demand Running Standalone On z/OS



Capacity Used High Priority Steady State - 85.2% CPU Minutes Unused (wasted) - 14.8% CPU Minutes **Priority Workload Metrics** Total Throughput: 417.8K Maximum TPS 129.7

Priority Transactional Workload On z/OS Does Not Degrade When Low Priority Donor Workload Is Added

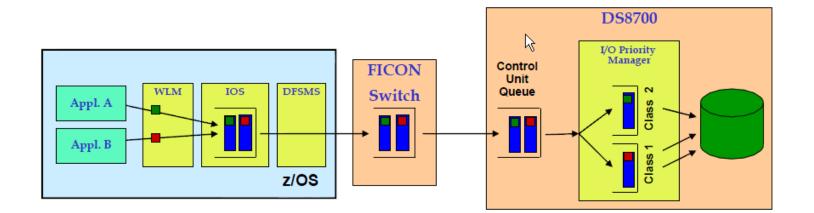


Unused (wasted) - 0% CPU Minutes

Total Throughput: 414.7K Maximum TPS 128.1

z/OS Workload Management Extends Priority All The Way Down To Storage

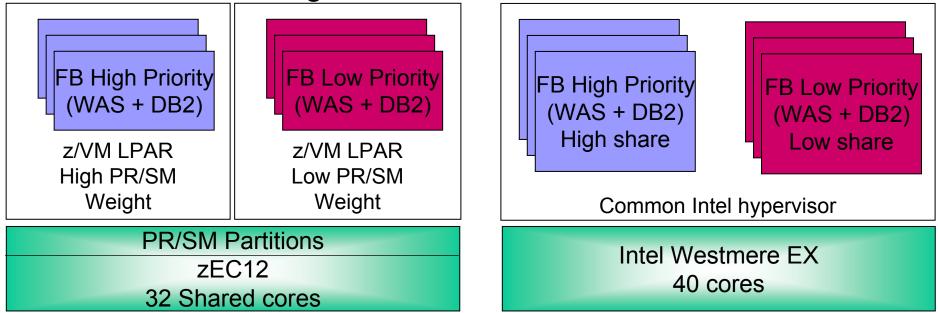
- FICON protocol supports advanced storage connectivity features not found in x86
- Priority Queuing:
 - Priority of the low-priority programs will be increased to prevent high-priority channel programs from dominating lower priority ones



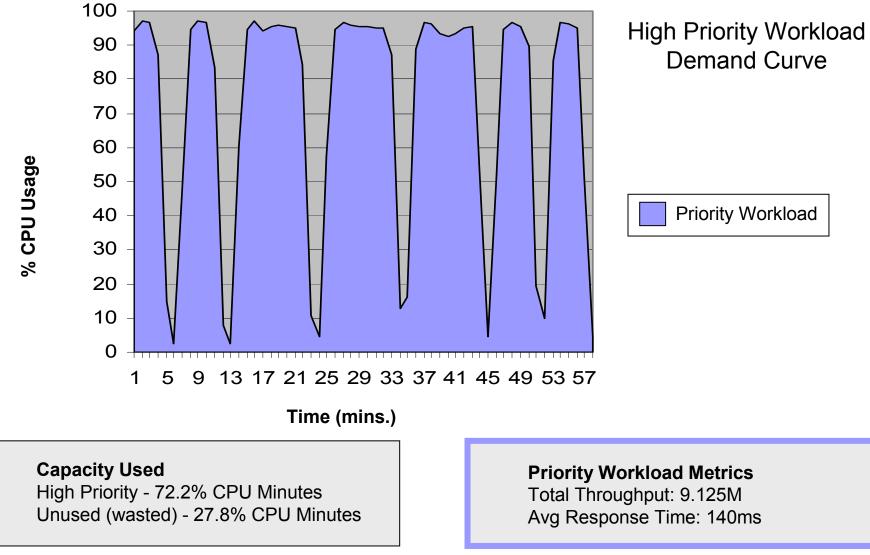
Intel can't do this

Comparison of System z To Intel Common Virtualization Environments

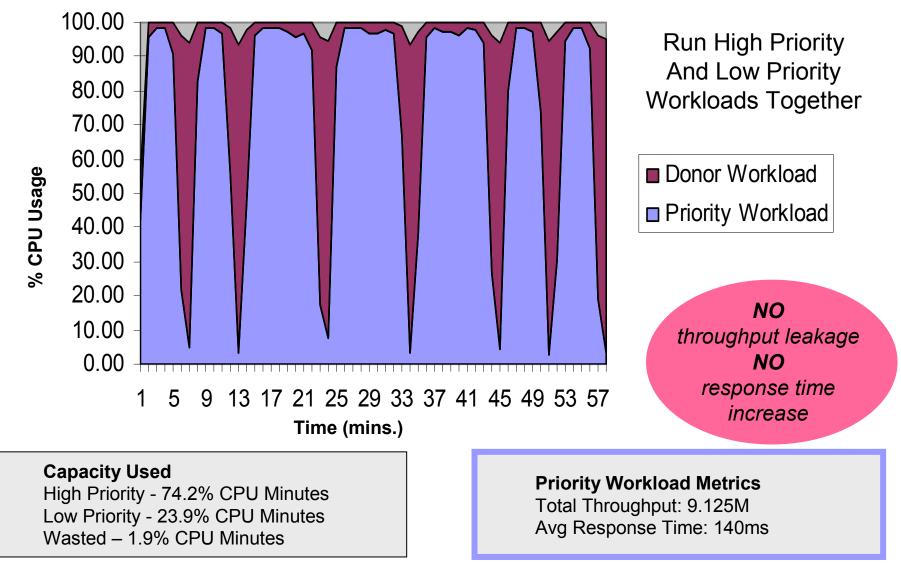
- High Priority web workload has defined demand over time
- SLA requires that response time does not degrade
- Low Priority web workload has unlimited demand
- It "soaks up" unused CPU minutes



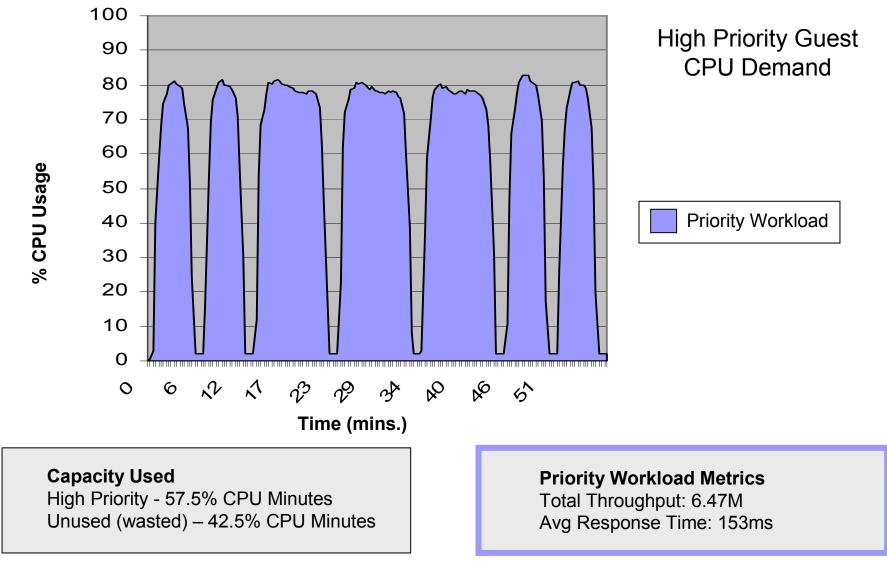
Priority Workload With Varying Demand Running Standalone On System z PR/SM



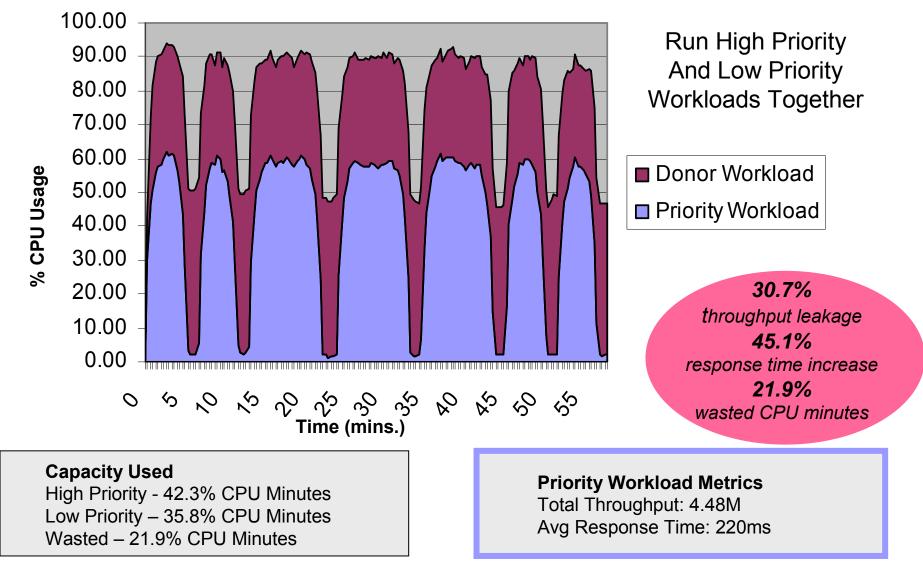
Priority Workload On System z Does Not Degrade When Low Priority Donor Workload Is Added



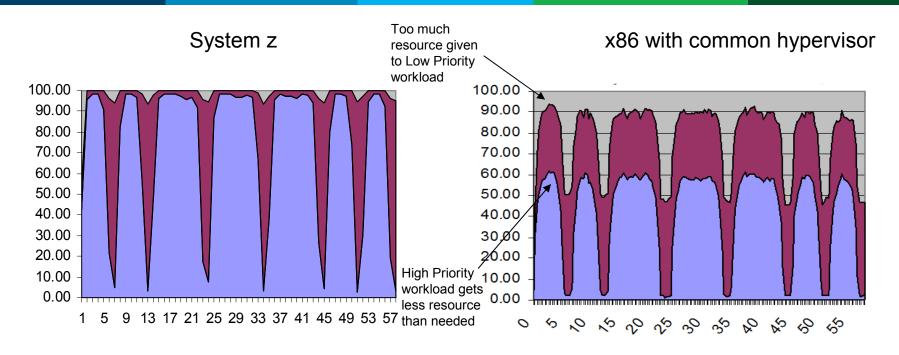
Priority Workload With Varying Demand Running Standalone On x86 Hypervisor



Priority Workload On x86 Hypervisor Degrades Severely When Low Priority Workload Is Added



System z Virtualization Enables Mixing Of High And Low Priority Workloads Without Penalty



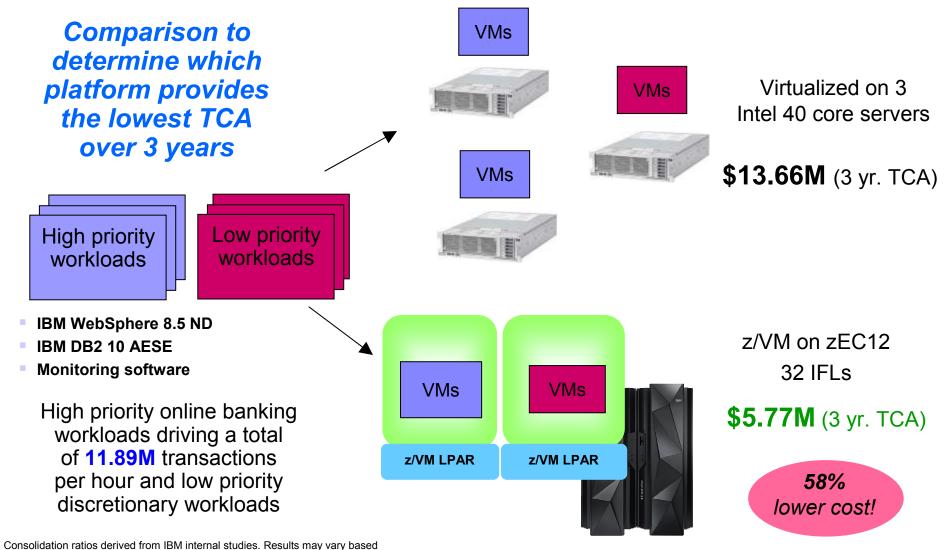
- Perfect workload management
- Consolidate workloads of different priorities on the same platform
- Full use of available processing resource (high utilization)

- Imperfect workload management
- Forces workloads to be segregated on different servers
- More servers are required (low utilization)

Inefficient Workload Management has an Economic Impact

- With Intel virtualization, the only practical solution to maintain SLA of high priority workload is to deploy workloads into separate environments
 - Most Intel virtualized deployments separate Dev environments from Production environments
- The need to maintain multiple environments directly affects total cost
 - Fixed size Intel boxes can force additional boxes to accommodate "spill over" high priority work
 - Spare capacity on additional machine is wasted as nothing else can run on it without impacting primary workload SLA
 - Additional environment needed to deliver lower priority workloads

Deliver High And Low Priority Workloads Together While Maintaining Response Time SLA



on customer workload profiles/characteristics. Prices will vary by country.

What System z Can Do That Intel-based Systems Can't

Workload Management Summary

- Additional environments needed for Intel can have a cost impact beyond server (hw & sw) acquisition
 - Supporting infrastructure components contribute to additional cost (TCO)
 - Storage, data copies, network, labor, environmental
- Imperfect workload management is one of the factors that leads to core proliferation in an offload scenario
 - Different environments that co-existed on z end up taking dedicated resources, driving up the total number of cores needed
- The perfect workload management capability of System z makes it an ideal Private Cloud platform
 - Support multiple tenants with different priorities
 - Ability to maintain priority enables maximum utilization of the platform, driving down cost per workload

What System z Can Do That Intel-Based Systems Can't

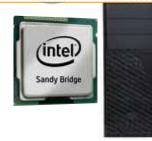


Servers Based on latest Intel technology (Sandy Bridge)

System z

Why Core Proliferation Happens When Moving Workload From System z To Intel

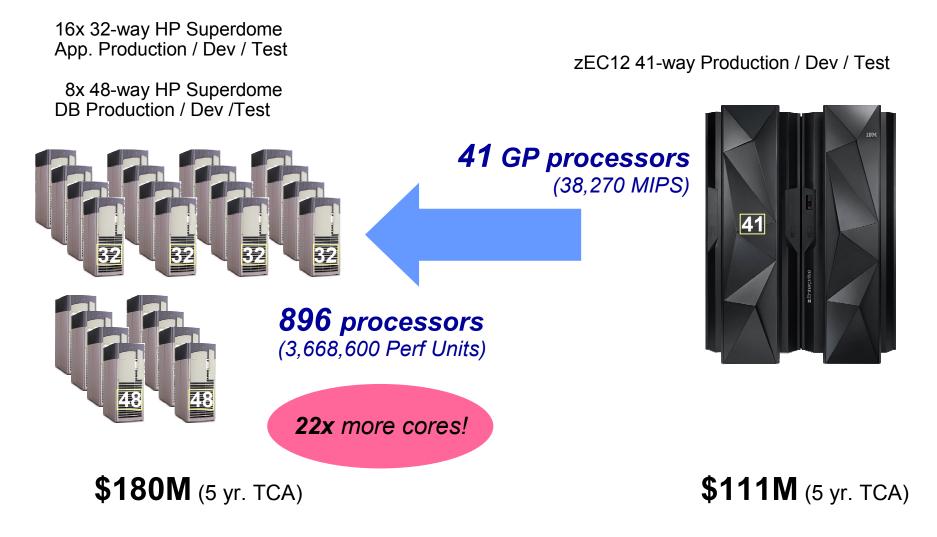
- De-consolidation of applications to dedicated servers – decomposing highly tuned co-located components
- Processing expansion requirements for CICS/COBOL applications
- 3x expansion when converting hierarchical databases to relational
- Functional segregation into production, development and test
- 100% hardware coverage for Disaster Recovery costs double





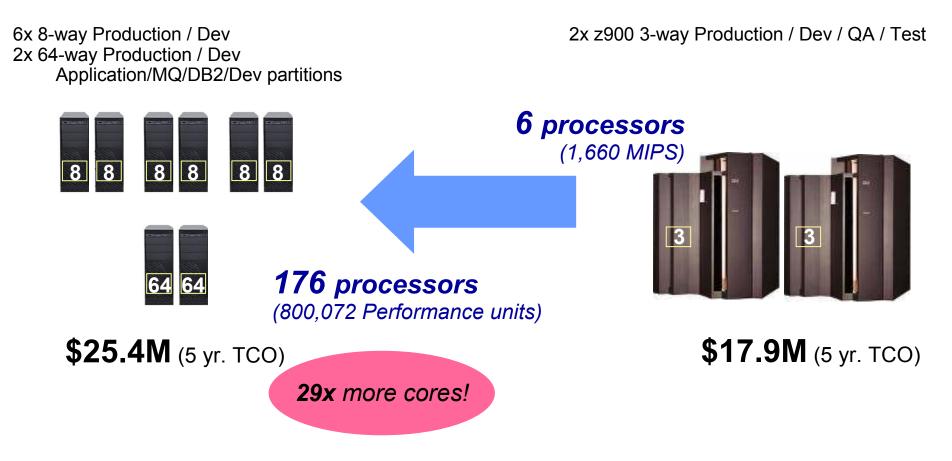


Core Proliferation For A Large Workload



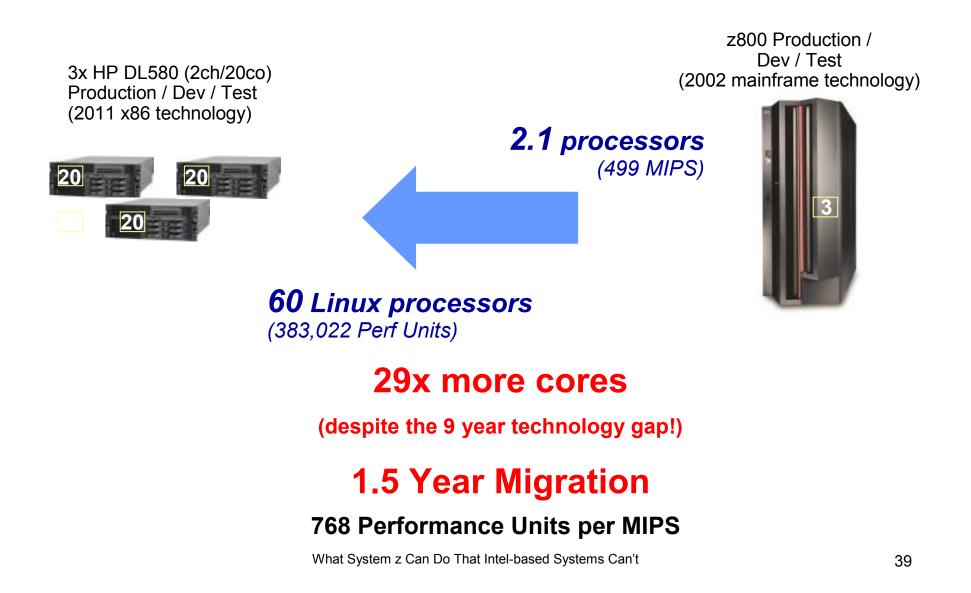
NOTE: To cover DEV/QA capacity, add 100% servers for distributed servers, add 25% MIPS (8,000) to System z

Core Proliferation For A Mid-sized Workload



482 Performance Units per MIPS

Core Proliferation For a Small Workload



Core Proliferation For Oracle Workloads

TCO study for a Media and Entertainment Industry customer



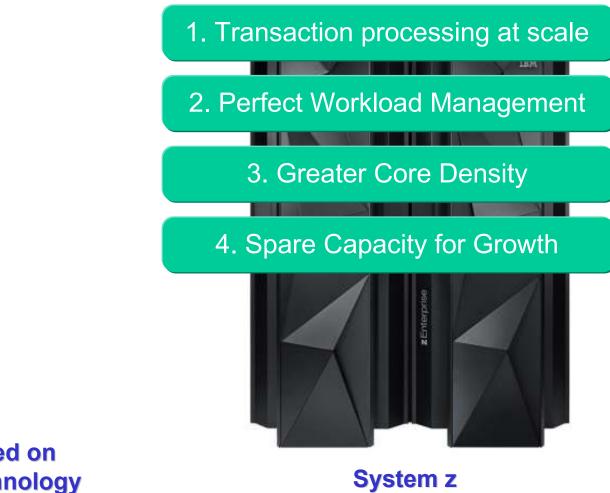
Hardware	\$2.9M
Software	\$24.2M
Labor	\$7.9M
Space, Power and cooling	\$1.2M
Disaster Recovery	\$6.5M
Total (5 yr. TCO)	\$42.7M



Disaster recovery Total (5 yr. TCO)	\$4.8M \$20.5M
Space, Power and cooling	\$0.5M
Labor	\$1.8M
Software	\$8.5M
Hardware	\$4.9M

Intel: Oracle DB + App costs = \$13.1M (LIC + maint over 5 yrs.). IBM: Oracle DB + App costs = \$1.92M (LIC + maint over 5 yrs.)

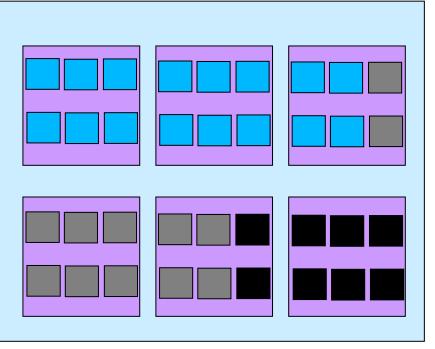
What System z Can Do That Intel-Based Systems Can't



Servers Based on latest Intel technology (Sandy Bridge)

System z Capacity On Demand Provides Elasticity To Handle Unexpected Peaks

- Capacity on Demand
 - "Books" are shipped fully populated
 - Activate dormant processors as needed
 - Use for temporary or permanent capacity
 - Self-managed on/off
- New capacity is immediately available for work without service disruption





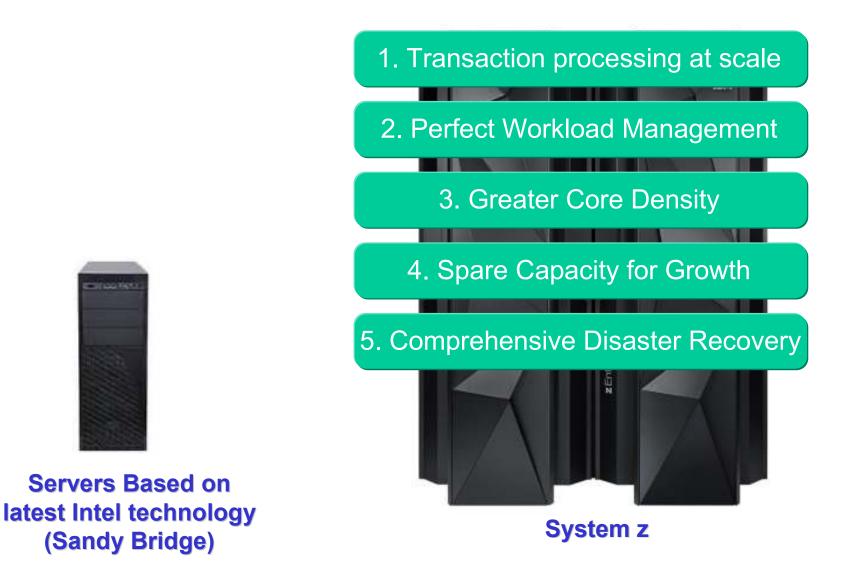
Active processors – pay full price

Inactive processors (On/Off CoD) – pay only 2% of full price

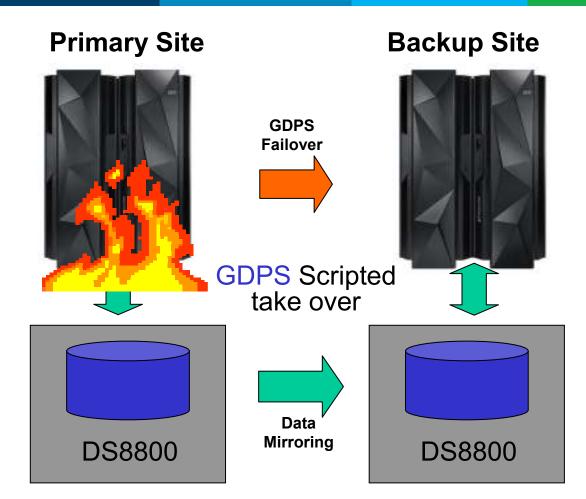
Dark processors (unused) - no charge

One Book with 36 Processors

What System z Can Do That Intel-Based Systems Can't



System z Disaster Recovery Is Systematic And Comprehensive



Site Failover

Failover to secondary site in case of complete site failure

```
Data Mirroring
```

 Protect data in the event of a disk system failure

Supports systematic Disaster Recovery for virtualized Linux environments also

Complexity Of Intel Disaster Recovery Solutions Prohibits Wide Spread Use

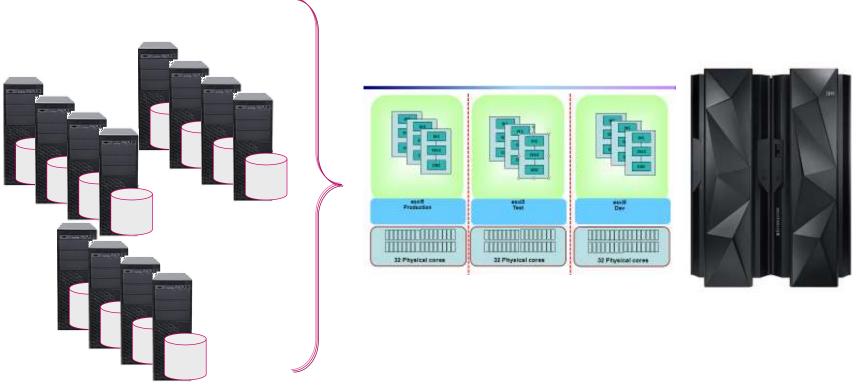
- Workloads on standalone
 Intel servers require a disaster recovery solution
 for each server
 - Data mirroring
 - ► Failover

- Embedded storage is difficult to mirror
- Comprehensive workload failover is not feasible for hundreds of servers



Consolidation Of Workloads On System z Simplifies Disaster Recovery

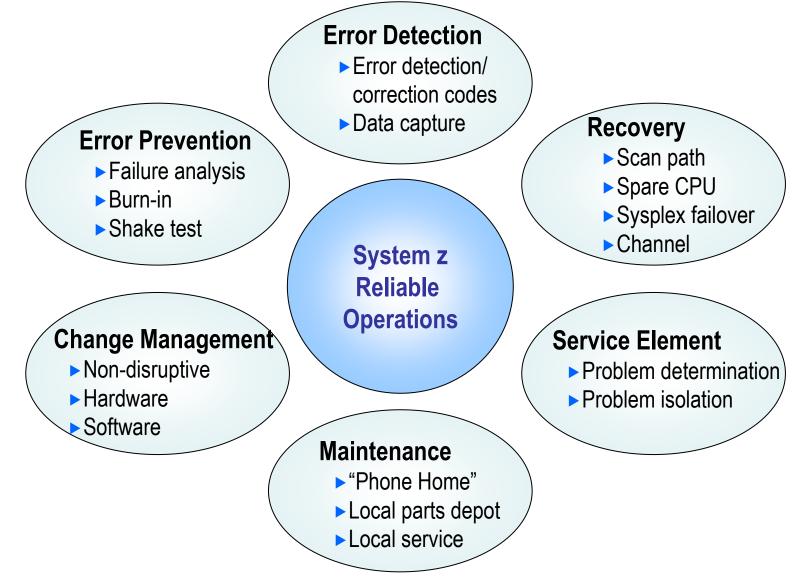
- Workloads are consolidated onto z/VM partitions as Linux guests
- Linux on System z can be failed over as part of GDPS



What System z Can Do That Intel-Based Systems Can't

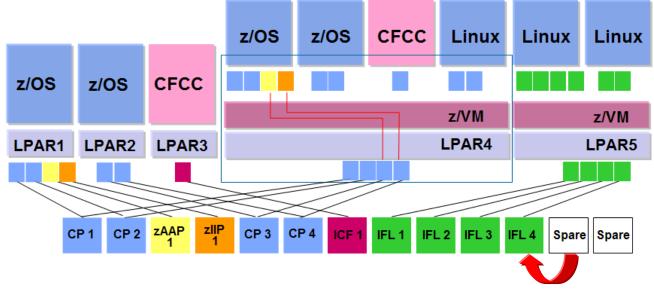


System z Has More Comprehensive Protection To Ensure Better Availability Than Intel

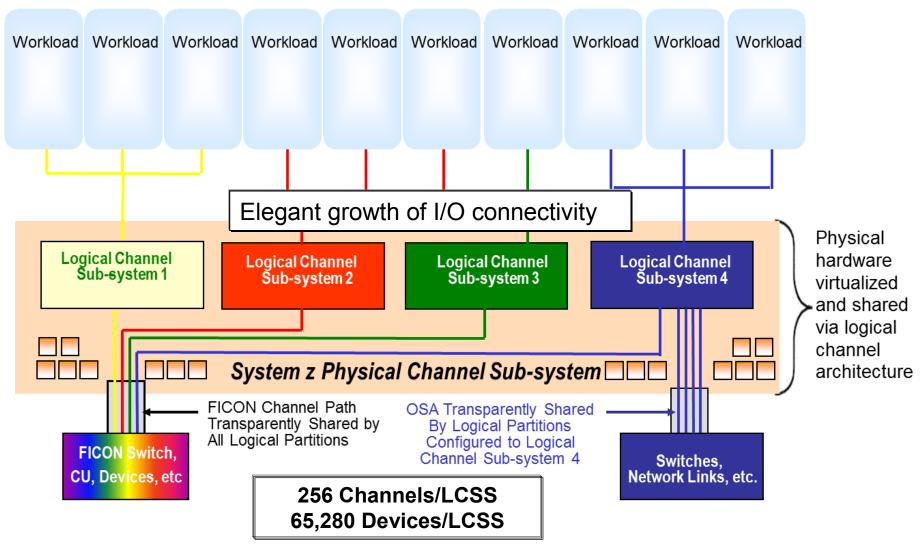


Example: CPU Sparing

- zEC12 has 2 spare CPUs per server
- System controllers can detect a failing processor chip
- Status of the unit of workload running on the failing CPU can be saved
- Failing CPU can be switched with the spare with NO interruption with the workload
- Alternatively, spare processors can be enabled at certain times during unexpected peak workloads
 - Another aspect of Capacity on Demand (COD)



Example: I/O Channel Failover



What System z Can Do That Intel-based Systems Can't

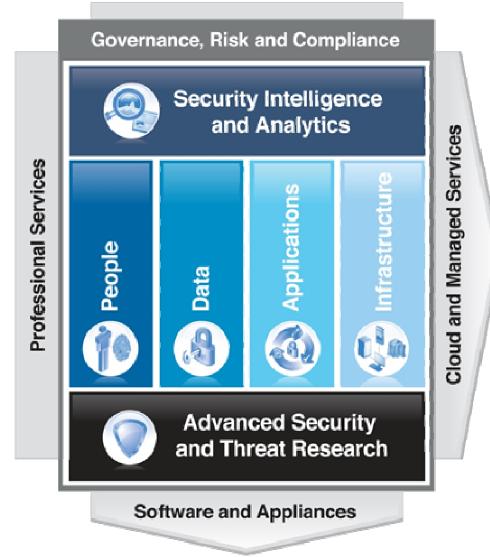
System z Supports Concurrent Operations During Hardware Repair – Intel Can't

Capability	zEC12	x86	
ECC on Memory Control Circuitry	Transparent While Running	Can recognize/repair soft errors while running limited ability with hard errors	
Oscillator Failure	Transparent While Running Must bring server down to replace		
Core Sparing	Transparent While Running	Must bring server down to replace	
Microcode Driver Updates	While Running	Some OS-level drivers can update while running, not firmware drivers; reboot often required	
Book Additions, Replacement	While Running Must bring server down to replace memory controllers, cache, et		
Memory Replacement	While Running	Must bring server down to replace	
Memory Bus Adaptor Replacement	While Running	Must bring server down to replace	
I/O Upgrades	While Running	Must bring server down to replace (limited ability to replace I/O in some servers)	
Concurrent Driver Maintenance	While Running	Limited – some drivers replaceable while running	
Redundant Service Element	2 per System	"Support processors" can act as poor man's SE, but no redundancy	

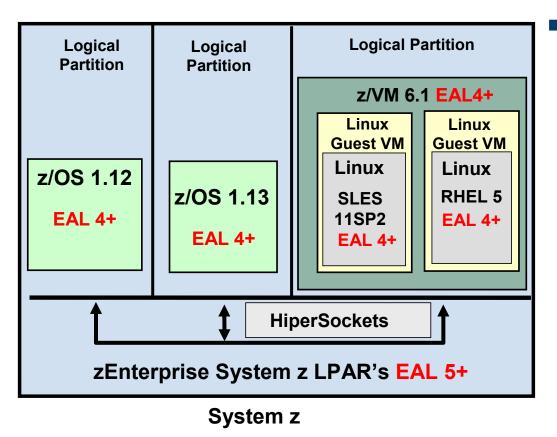
What System z Can Do That Intel-Based Systems Can't



Enterprise Security Requires Many Elements: Infrastructure



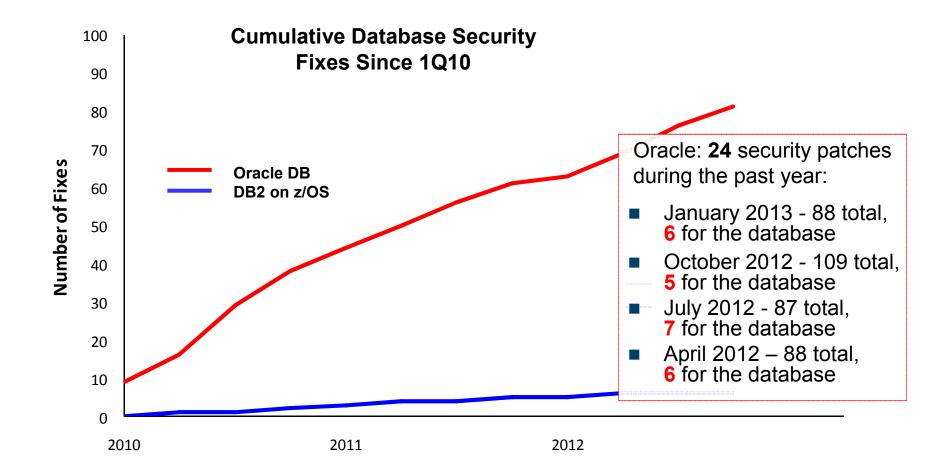
Common Criteria Certifications Show System z Platform Security Leads the Industry



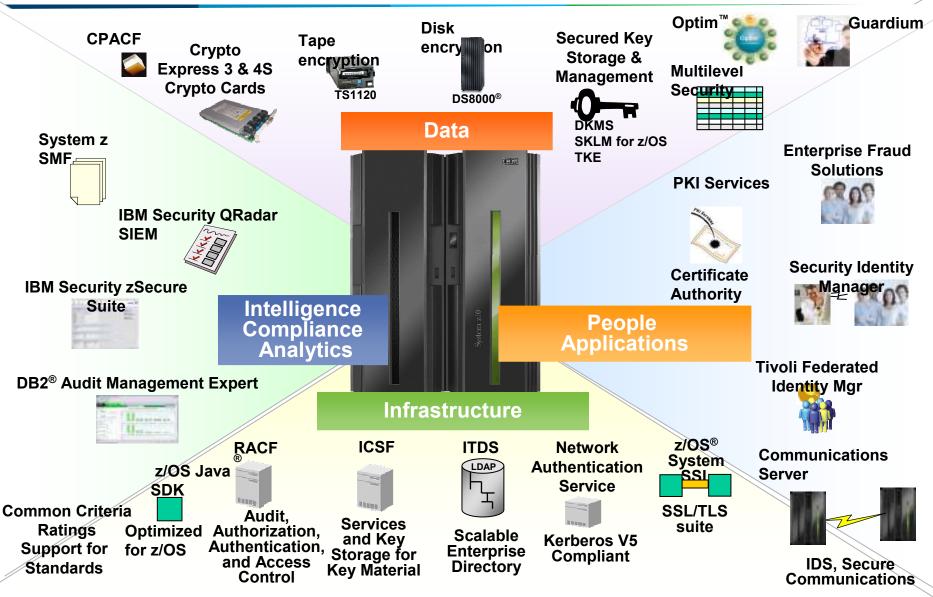
- What is Common Criteria?
 - Common Criteria is an accepted standard for evaluating the inherent security of a computing system
 - Common Criteria is based on a set of functional and assurance requirements
 - A higher Enterprise Assurance Level (EAL) rating is more secure
 - The security requirements in Common Criteria have gained support as "best practices"
 - IBM System z holds the highest EAL grades in Common Criteria!

EAL Ratings: https://www.bsi.bund.de/SharedDocs/Zertifikate/CC/Betriebssysteme/0852.html

DB2 Maturity Delivers A Proven Track Record For Data Security



Elements Of System z Security



Virtualized System z security is superior to other platforms and augmentation costs less

Security Level Description	MF	x86	UNIX
Normal corporate	100.00%	21.00%	35.00%
Credit card processing involved	100.00%	14.00%	26.00%
Banking	100.00%	8.00%	14.00%
Healthcare	100.00%	7.50%	11.00%
Research	78.00%	3.00%	8.00%
Defense	64.00%	1.00%	3.00%

Security Natively Covered by Platform

Major security deficiencies on distributed platforms

Distributed platforms require *considerable additional expense*

On System z most security requirements are standard

Little additional augmentation required on System z

Security Level Description	MF	x 86	UNIX
Normal corporate	0.00%	25.20%	12.10%
Credit card processing involved	0.00%	38.40%	16.90%
Banking	0.00%	63.70%	22.40%
Healthcare	0.00%	81.60%	30.70%
Research	2.10%	134.80%	56.90%
Defense	4.30%	187.90%	97.50%

Incremental Cost to Achieve Required Security

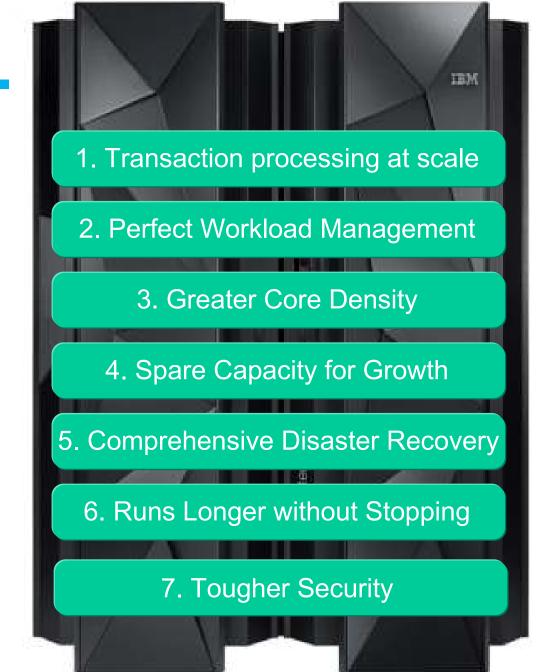
Source: "Comparing Virtualization Alternatives – What's best for your business?"

© 2012, Solitaire Interglobal Ltd. http://public.dhe.ibm.com/common/ssi/ecm/en/zsl03192usen/ZSL03192USEN.PDF

The Choice Is Clear!

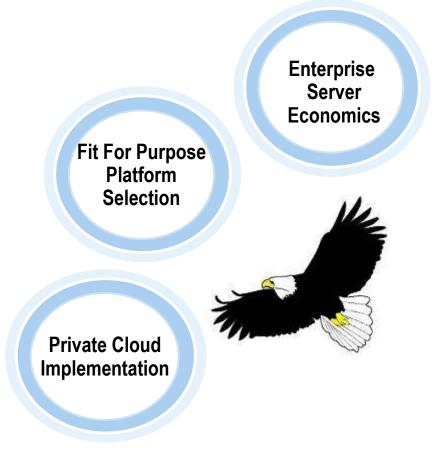
System z is better than Intelbased systems for Systems of Record





The IBM Eagle team helps customers understand mainframe costs and value

- Worldwide team of senior technical IT staff
- Free of Charge Total Cost of Ownership (TCO) studies
 - Help customers evaluate the lowest cost option among alternative approaches
 - Includes a one day on-site visit and is specifically tailored to a customer's enterprise
- Studies cover POWER, PureSystems and Storage accounts in addition to System z
 - For both IBM customer and Business Partner customer accounts
- Over 300 customer studies since formation in 2007
- Contact: eagletco@us.ibm.com



IEM 👸

59

For More Information please contact...

Len Santalucia, CTO & Business Development Manager Vicom Infinity, Inc. One Penn Plaza – Suite 2010 New York, NY 10119 212-799-9375 office 917-856-4493 mobile lsantalucia@vicominfinity.com

About Vicom Infinity

Account Presence Since Late 1990's **IBM Premier Business Partner** Reseller of IBM Hardware, Software, and Maintenance Vendor Source for the Last 4 Generations of Mainframes/IBM Storage Professional and IT Architectural Services Vicom Family of Companies Also Offer Leasing & Financing, Computer Services, and IT Staffing & IT Project Management What System z Can Do That Intel-based Systems Can't